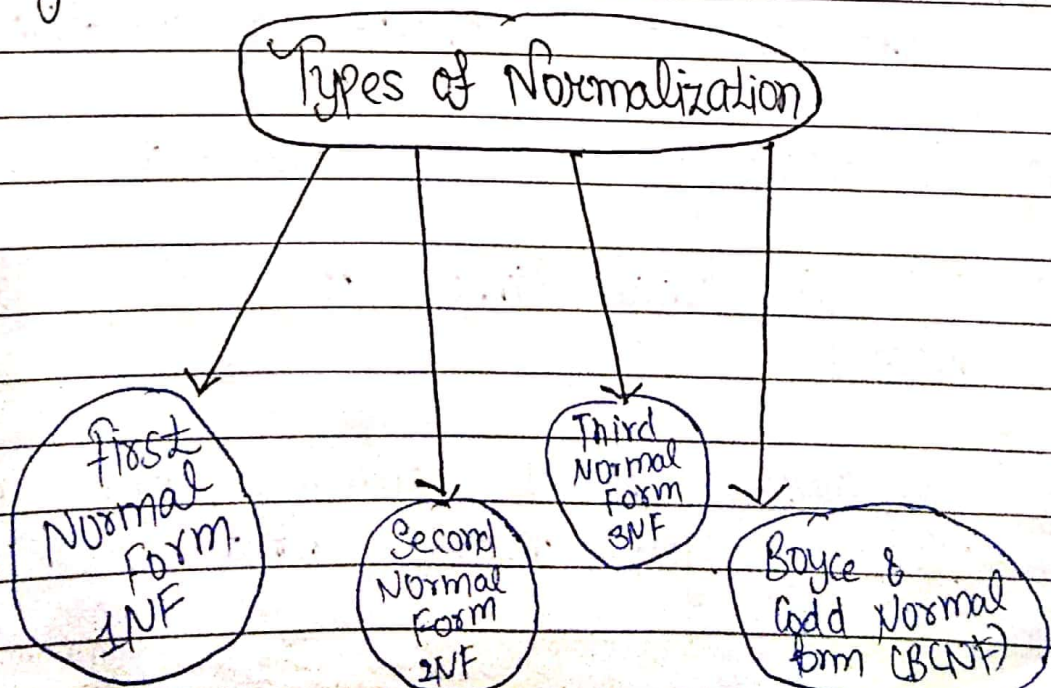


# #DBMS Notes.

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## \* Normalization :-

- Normalization is the process of organizing the data in the database.
- Normalization is used to minimize the redundancy from a relation or set of relations.
- Normalization is based on the concept of Normal forms.
- Normal forms are used to eliminate or reduce redundancy in database tables.
- It is used to avoid the data redundancy, insertion anomaly, update anomaly & delete anomaly.



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Five Normal forms These are 1NF, 2NF, 3NF, 4NF, 5NF where NF stands for Normal forms.

→ The 1NF, 2NF, 3NF were proposed by Dr. E.F. Codd. Later on, 3NF known as BCNF was proposed by Boyce and Dr. Codd.

→ The 4NF and 5NF Normal forms were proposed after BCNF.

\* Advantages of Normalization :-

- (i) Normalization helps to minimize the data redundancy.
- (ii) Better understanding of data.
- (iii) More efficient data structure.
- (iv) More flexible data structure.
- (v) Enforces the concept of Relational integrity.

\* Disadvantages of Normalization :-

(i) It is very time consuming and difficult process.

(ii) you cannot start building the database before you know what the user needs.

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(c) First Normal Form (1NF) :-

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- All Relation is said to be in 1NF if and only if it follows the Rules :-
  - A relation will be 1NF if it contains an atomic value.
  - There are no repeating groups in the table.
  - Table should not contain any multivalued attribute.
- It is a technique to Remove or Reduce Redundancy from a table.
- It must hold only single-valued attribute.

Example :- Relation EMPLOYEE is not in 1NF because of multi-valued attributes

EMP ID	EMP_NAME	EMP_PHONE
14	John	728226365 90643828
20	Harry	7393702
12	Sam	858903022 85747838

→ The decomposition of the EMPLOYEE table into 1NF :-

EMP ID	EMP_NAME	EMP_PHONE
14	John	728226365
14	John	90643828
20	Harry	7393702
12	Sam	858903022
12	Sam	85747838

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## (a) Second Normal Form :-

- In the 2NF, relational must be in 1NF.
- In the second normal form, all non-key attributes are fully functional dependent on the primary key.

Example: A school can store the data of teachers and the subjects they teach. In a school, a teacher can teach more than one subject.

Teacher table

TEACHER_ID	SUBJECT	TEACHER_AGE
25	Java	30
25	Python	30
47	DBMS	35
83	Maths	38
83	Computer	38

- non-prime attribute is TEACHER\_AGE is dependent on TEACHER\_ID.
- To convert the given table into 2NF, we decompose it into two tables.

TEACHER ID	TEACHER AGE
25	30
47	35
83	38

TEACHER ID	SUBJECT
25	Java
25	Python
47	DBMS
83	Maths
83	Computer

- To be in second normal form, a relation must be in 1NF and relation must not contain any partial dependency

Partial Dependency :- A Partial dependency would occur whenever a non-prime attribute depends functionally on part of candidate key.

The 2NF (Second Normal form) eliminates the partial dependency.

### (3) Third Normal form (3NF) :-

- o 3NF is an upgrade to second normal form.
- o When a table is in the Second Normal form and has no transitive dependency, then it is the Third Normal Form.

Transitive Dependency :- When an indirect relationship causes functional dependency it is called Transitive Dependency.

Example: If  $P \rightarrow Q$  and  $Q \rightarrow R$  is true then  $P \rightarrow R$  is a transitive dependency.

#### Rules for 3NF :-

(i) Table should be in 2NF.

(ii) There is no transitive dependency.

Example: Consider a relation student (rollno, game, fee)

Rollno	Game	fee
1	Basketball	500
2	Basketball	500
3	Basketball	500
4	Cricket	600
5	Cricket	600
6	Tennis	400

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F - (rollno, → game, rollno → fee, game → fee).

Rollno<sup>+</sup> = (rollno, game, fee)

⇒ rollno is primary key.

→ fee has transitive dependency on rollno via game.

### Anomalies

(i) Insertion anomaly :- A new game can't be inserted into the table unless we get a student to play that game.

(ii) Deletion anomaly :- If rollno G is deleted from the table we also lost the complete information regarding tennis.

(iii) Update anomaly :- To change the fee for basketball we need to make changes in more than one place.

### Decomposition for 3NF :-

To overcome these anomalies, the student table should be divided into smaller tables.

So divide the student table into R1 (game, fee) and R2 (rollno, game).

Q1

Rollno	Game
1	Basketball
2	Basketball
3	Basketball
4	Cricket
5	Cricket
6	tennis

Q2

Game	fee
Basketball	500
Cricket	600
Tennis	400

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## \* Boyce Codd Normal Form (BCNF) :-

- o BCNF is the advance version of 3NF. It is stricter than 3NF.
- o A relation in BCNF, if it is already in 3NF and determinant is a candidate key.
- o Determinant is a simple attribute or composite attributes which can uniquely determines the value of other attributes.

Example - Table : ITEM

Item No.	Item-Name	Quantity	Price
1001	Pen	20	50
1001	Pencil	15	10
1004	Eraser	30	10
1007	Pen	25	50

→ In this, attribute Quantity is fully functionally dependent on (Item No, item Name)

(Item No, item Name) → Quantity

Composite determinant

→ The attribute Price is functionally dependent on item name

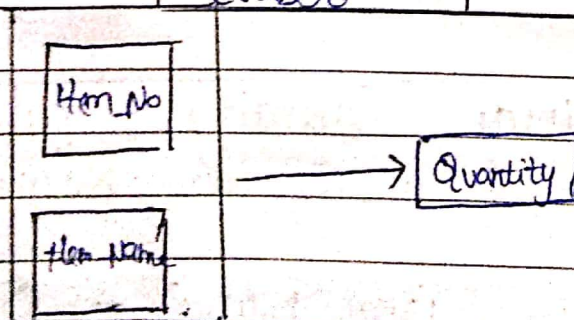
item Name  $\Rightarrow$  Price

$\hookrightarrow$  simple attribute determinant.

→ The normalized relations / tables for above dependencies are shown below:-

item No	item-Name	Quantity
1001	Pen	20
1001	Pencil	15
1004	Eraser	30
1007	Pen	25

Item Name	Price
Pen	50
Pencil	10
Eraser	10



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## \* Fourth Normal Form (4NF) :-

A Relation 'R' is in fourth Normal form (4NF) if it is in BCNF and has no multi value dependency.

Multi valued dependency :- A table is said to have multi-valued dependency, if the following conditions are true,

1. for a dependency  $A \twoheadrightarrow B$ , if for a single value of A, multiple value of B exists, then the table may have multi-valued dependency.
2. A table should have atleast 3 Columns for it to have a multi-valued dependency.
3. for a relation  $R(A, B, C)$  if there is a multi-valued dependency between A and B, then B and C should be independent of each other.

Example :-

$A \twoheadrightarrow B$

and

$A \twoheadrightarrow C$

Name	department	Language
John	Sales	English
John	Sales	French
John	Marketing	English
Michael	Sales	English
Michael	Sales	French
Michael	Sales	Japanese

In this Table:-

The attribute Department is multi valued dependent on attribute Name.

Name  $\twoheadrightarrow$  Department

The attribute language is multivalued dependent on attribute Name.

Name  $\twoheadrightarrow$  language.

→ So convert the above relation (EMP) with <sup>multi-</sup>valued dependencies into 4NF, move each multi-valued dependency pair to a new table.

EMP

Name	Department	Language
John	Sales	English
John	Sales	French
John	Marketing	English
Michael	Sales	English
Michael	Sales	French
Michael	Sales	Japanese

Table EMP-DEPT

Name	Department
John	Sales
John	Marketing
Michael	Sales

Table EMP-LANG

Name	Language
John	English
John	French
Michael	English
Michael	French
Michael	Japanese

(5) Fifth Normal form (5NF):-  
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- A relation R in the fifth Normal form (5NF) if it is already in 4NF and has no join dependency and joining should be lossless.
- 5NF is satisfied when all the tables are broken into as many tables as possible in order to avoid redundancy.
- 5NF is also known as Project-Join Normal form (PJ/NF).

Example:-

SUBJECT	LECTURER	SEMESTER
Computer	Anshika	Sem 1
Computer	John	Sem 1
Math	John	Sem 1
Math	Palvi	Sem 2
DBMS	Ram	Sem 1

So to make the above table into 5NF, we can decompose it into three Relations P1, P2 & P3.

P1

Semester	Subject
Sem 1	Computer
Sem 1	math
Sem 1	DBMS
Sem 2	math

P2

Subject	Lecturer
Computer	Anshika
Computer	John
math	John
math	Palvi
DBMS	Ram

P3

Semester	Lecturer
Sem 1	Anshika
Sem 1	John
Sem 1	John
Sem 2	<del>Anshika</del> Palvi
Sem 1	Ram